Efficient Low-Complexity Antenna Selection Algorithms in Multi-User Massive MIMO Systems with Matched Filter Precoding

Zaid Abdullah, Member, IEEE, Charalampos C. Tsimenidis, Senior Member, IEEE, Gaojie Chen, Senior Member, IEEE, Martin Johnston, Member, IEEE, and Jonathon A. Chambers, Fellow, IEEE

Abstract—In this work, two efficient low complexity Antenna Selection (AS) algorithms are proposed for downlink Multi-User (MU) Massive Multiple-Input Multiple-Output (M-MIMO) systems with Matched Filter (MF) precoding. Both algorithms avoid vector multiplications during the iterative selection procedure to reduce complexity. Considering a system with \(N\) antennas at the Base Station (BS) serving \(K\) single-antenna users in the same time-frequency resources, the first algorithm divides the available antennas into \(K\) groups, with the \(k\)th group containing \(N/K\) antennas that have the maximum channel norms for the \(k\)th user. Therefore, the Signal-to-Interference plus Noise Ratio (SINR) for the \(k\)th user can be maximized by selecting a subset of the antennas from only the \(k\)th group, thereby resulting in a search space reduction by a factor of \(K\). The second algorithm is a semi-blind interference rejection method that relies only on the signs of the interference terms, and at each iteration the antenna that rejects the maximum number of interference terms will be selected. The performance of our proposed methods is evaluated under perfect and imperfect Channel State Information (CSI) and compared with other low complexity AS schemes in terms of the achievable sum rate as well as the energy efficiency. In particular, when the Signal-to-Noise Ratio (SNR) is 10 dB, and for a system with 20 MHz of bandwidth, the proposed methods outperform the case where all the antennas are employed by 108.8 and 49.2 Mbps for the first and second proposed algorithms, respectively, given that the BS has perfect CSI knowledge and is equipped with 256 antennas, out of which 64 are selected to serve 8 single-antenna users.

Index Terms—Massive MIMO, Antenna Selection, Multi-User, Complexity Reduction, Sum Rate, Energy Efficiency

I. INTRODUCTION

MASSIVE Multiple-Input Multiple-Output (M-MIMO) has been a hot topic of research in wireless communications in recent years, and it refers to a system where tens to hundreds of antennas are placed at the Base Station (BS) to serve a much smaller number of users. Compared to conventional MIMO systems, M-MIMO shows a great improvement in capacity and reliability while reducing the radiated energy dramatically [1], since increasing the number of antennas will provide additional degrees of freedom, which in return can be utilized to enhance the energy consumption of such systems [2]. For downlink transmission, precoding techniques should be applied to map the \(K \times 1\) symbol vector intended for the \(K\) single-antenna users into an \(N \times 1\) data vector before being transmitted through the wireless channel, where \(N\) is the number of antennas at the BS, \(N \gg K\). The two most common types of linear precoding are Matched Filter (MF) and Zero Forcing (ZF) [3]. However, in this paper we consider only MF precoding since ZF does not only require higher complexity for finding the pseudo inverse of the channel matrix, but also can suffer from serious performance limitations. For example, one of the main limitations of ZF is that when the number of users grows large for a fixed number of antennas at the BS, the ZF precoder suffers from a sum rate loss [4], since most of the power will be used to null the interference. Moreover, the ZF suffers from degraded performance when there are users located on the cell-edge, also known as the “near-far” problem. In addition, the ZF precoder is very sensitive to different types of distortions such as unmodeled interference [5], [6]. In contrast, MF precoding requires lower processing complexity and can accommodate more users than ZF.

Although increasing the number of antennas at the BS offers great advantages in terms of performance, it does come at the price of increased hardware complexity, cost, and power consumption. Employing all the antennas means that each antenna should be connected to a separate Radio Frequency (RF) chain, and each RF chain consists of a mixer, analogue to digital converter, and amplifier. Therefore, and unlike the antenna elements, RF chains are expensive. Moreover, RF chains are considered as highly power demanding elements, and they consume 50%-80% of the total transceiving power [7]. One way to reduce the hardware complexity, cost, and power consumption while keeping the advantages of a M-MIMO system, is by applying Antenna Selection (AS) schemes [8]. Therefore, designing an efficient low complexity AS algorithm is an important topic in M-MIMO systems.

There has been a considerable amount of work on AS with both conventional and M-MIMO systems, for example, in [4], it was found that not all antennas contribute equally in M-MIMO systems, and AS can reduce the complexity and cost of M-MIMO without large degradation in the performance; while the authors in [9] designed an AS scheme to reject
the co-channel interference for a single user uplink scenario. Moreover, the authors in [10] and [11] aimed to enhance the error performance through AS by exploiting the temporal correlation, and the constructive interference, respectively. Although the proposed algorithm in [11] has low complexity, it works efficiently only for low modulation Phase Shift Keying (PSK) signalling, and it is data dependent, which means that extremely fast RF switching is required. Furthermore, the authors in [12] and [13] addressed the design of AS algorithms to maximize the Energy Efficiency (EE) of point-to-point M-MIMO systems. The authors in [14] proposed a novel iterative AS method for an uplink point-to-point M-MIMO system with a Maximum Ratio Combining (MRC) receiver, under spatially correlated channels. Their work focused on minimizing the Mean Square Error (MSE) of the received signal to improve the error rate performance for a single user scenario. The authors in [15] proposed a Rectangular Maximum-Volume (RMV) theory based AS technique, for downlink M-MIMO systems. However, their work focused on maximizing the sum rate capacity, which can be obtained via high complexity dirty paper coding. Moreover, the authors in [16] and [17] designed an AS algorithm with quartic complexity in point-to-point M-MIMO systems, while the authors in [18] studied the trade-off between energy and spectral efficiencies under random AS in Multi-User (MU) M-MIMO downlink systems. Moreover, in [19], the authors proposed an AS scheme under interference alignment, where the interference is forced to zero through transmitter-receiver beamforming. The authors in [20] proposed an AS method for power minimization in MU downlink M-MIMO systems, under a predefined Quality of Service (QoS) requirement. In addition, the authors in [21] and [22] proposed a joint AS and user scheduling scheme with ZF precoding. Finally, the authors in [23] showed that when MF precoding is applied, removing certain antennas can improve the sum rate performance of the system, and they proposed a low complexity greedy AS algorithm for sum rate maximization in downlink M-MIMO systems.

In this paper, we design two novel AS algorithms for downlink MU M-MIMO systems for Signal-to-Interference plus Noise Ratio (SINR) maximization with low complexity. Our contributions in this paper, which to the best of our knowledge have not been considered in any previously published work, can be summarized as follows:

1) We design a User-Centric AS (UCAS) algorithm with reduced search space, where the available antennas are divided into \( K \) groups, and each group corresponds to only one user, where the \( k \)th group contains the antennas that have the maximum channel norms for the \( k \)th user. Therefore, to maximize the SINR for the \( k \)th user at a given iteration, the proposed algorithm selects an antenna from only the \( k \)th group. This reduces the search space by a factor of \( K \).

2) The second proposed algorithm, called Semiblind Interference Rejection AS (SIRAS), is designed to reject the highest number of interuser interference terms at any given iteration, based on the signs of only the interference terms.

3) Both algorithms avoid any vector multiplications during the iterative selection process by storing the multiplications of any two entries in the channel matrix before the iterative algorithms start. This results in dramatic complexity reduction in terms of number of floating-point operations (FLOPs) required for their implementations.

4) The performance of the proposed algorithms is evaluated in terms of the achievable sum rate, energy efficiency, as well as computational complexity, and compared with other low complexity AS algorithms. Our proposed methods demonstrate an improved performance-complexity trade-off.

The rest of this paper is organized as follows. In Section II, we introduce the system model and formulate the AS problem. In Section III, the proposed AS algorithms are introduced and explained in detail. In Section IV, different numerical results are presented along with their discussions. The complexities of the different AS algorithms used in this work are evaluated in Section V in terms of number of FLOPs required for their implementations. Finally, the conclusions are drawn in Section VI.

### TABLE I

<table>
<thead>
<tr>
<th>Notation</th>
<th>Explanation</th>
</tr>
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<tbody>
<tr>
<td>(</td>
<td></td>
</tr>
<tr>
<td>( a^* )</td>
<td>The element-wise conjugate of vector ( a )</td>
</tr>
<tr>
<td>( a^H )</td>
<td>The Hermitian transpose of vector ( a )</td>
</tr>
<tr>
<td>( \Re[a] )</td>
<td>The real part of complex number ( a )</td>
</tr>
<tr>
<td>( \Im[a] )</td>
<td>The imaginary part of complex number ( a )</td>
</tr>
<tr>
<td>( 0_{M \times N} )</td>
<td>The ( M ) by ( N ) zero matrix</td>
</tr>
<tr>
<td>( I_N )</td>
<td>The ( N ) by ( N ) identity matrix</td>
</tr>
<tr>
<td>( a_i )</td>
<td>The ( i )th row of matrix ( A )</td>
</tr>
<tr>
<td>( a_j^T )</td>
<td>The ( j )th column of matrix ( A )</td>
</tr>
<tr>
<td>( a_{i,j} )</td>
<td>The ( i )th element of the ( j )th column of ( A )</td>
</tr>
<tr>
<td>(</td>
<td>a</td>
</tr>
<tr>
<td>( a! )</td>
<td>The factorial operator</td>
</tr>
<tr>
<td>( \binom{n}{k} = \frac{n!}{k!(n-k)!} )</td>
<td>The binomial coefficient</td>
</tr>
</tbody>
</table>

The list of notations used throughout this work alongside their explanations are shown in Table I. We next introduce the model for the MU M-MIMO system.

### II. SYSTEM MODEL AND PROBLEM FORMULATION

#### A. Channel, signal, and noise models

We consider a single cell MU M-MIMO system operating in the downlink scenario as depicted in Fig. 1. Time Division
Duplex (TDD) transmission is assumed, where the users send orthogonal pilots to the BS to obtain Channel State Information (CSI). Let \( \mathbf{H} = [\mathbf{h}_1, \ldots, \mathbf{h}_K]^T \in \mathbb{C}^{K \times N} \) be the channel matrix, where \( \mathbf{h}_k \in \mathbb{C}^{N \times 1} \) includes the channel coefficients between the BS and user \( k \), with \( h_{k,n} \sim \mathcal{CN}(0, \sigma_n^2) \), and let \( \mathbf{w}_k \in \mathbb{C}^{N \times 1} \) denote the unit norm precoding vector for the \( k \)th user, and it can be given as

\[
\mathbf{w}_k = \frac{\mathbf{h}_k^*}{\|\mathbf{h}_k\|}.
\]

Assuming that the BS has perfect knowledge of \( \mathbf{H} \), and MF precoding is applied, the transmitted signal vector from the BS to all users can be given as

\[
\mathbf{s} = \sum_{k=1}^{K} \sqrt{p_k} \mathbf{w}_k \mathbf{x}_k,
\]

where \( \mathbf{x} = [x_1, \ldots, x_K]^T \) is the vector of information symbols intended for the \( K \) users, with \( E\{\mathbf{x}\mathbf{x}^H\} = \mathbf{I}_K \), \( p_k \) is the power allocated for the \( k \)th user, and it follows the constraint

\[
\sum_{i=1}^{K} p_i = P_T,
\]

where \( P_T \) is the total transmission power in Watts available at the BS for all users. Moreover, the received signal at the \( k \)th user can be given as

\[
y_k = \mathbf{h}_k^T \mathbf{s} + n_k
\]

\[
= \sqrt{p_k} \mathbf{h}_k^T \mathbf{w}_k \mathbf{x}_k + \sum_{i=1}^{K} \sqrt{p_i} \mathbf{h}_k^T \mathbf{w}_i \mathbf{x}_i + n_k,
\]

where \( n_k \sim \mathcal{CN}(0, \sigma_n^2) \) is the Additive White Gaussian Noise (AWGN) at the \( k \)th user.

**B. SINR, achievable rates, and energy efficiency**

In this work, we will use the sum rate and EE metrics to demonstrate the efficiency of the proposed AS algorithms. For systems with MF precoding, the SINR for the \( k \)th user can be expressed as [24]

\[
\gamma_k = \frac{p_k |\mathbf{h}_k^T \mathbf{w}_k|^2}{\sum_{i=1, i \neq k}^{K} p_i |\mathbf{h}_k^T \mathbf{w}_i|^2 + \sigma_n^2},
\]

and the achievable rate for the same user is a function of the SINR, and it can be given as

\[
R_k = \log_2(1 + \gamma_k),
\]

while the total sum rate is the sum of achievable rates for all users, i.e.

\[
R = \sum_{k=1}^{K} R_k.
\]

Furthermore, the EE in bits/Joule can be defined as the total bandwidth, multiplied by the total sum rate over the total power consumed [13], i.e.

\[
\text{EE} = \frac{B \cdot R}{P_{\text{total}}},
\]

where \( B \) is the bandwidth, and \( P_{\text{total}} \) is the total power consumed at the transmitter and the receiver, and can be given as [25]

\[
P_{\text{total}} = \frac{P_T}{\eta} + N_{RF} P_{tx} + K P_{rx},
\]

where \( \eta \) is the power amplifier efficiency, \( N_{RF} \) is the number of activated RF chains at the transmitter, while \( P_{tx} \) and \( P_{rx} \) are the circuit power consumption per RF chain at the transmitter and the receiver, respectively [25].

**C. AS problem formulation**

For a BS equipped with only \( N_s \) RF chains, AS schemes can be applied to optimize a certain cost function. In this work, we aim to select \( N_s \) out of the available \( N \) antennas at the BS to maximize the SINRs. This optimization problem can be formulated as follows

\[
\text{maximize } \Delta \sum_{k=1}^{K} \left( \frac{p_k |\mathbf{h}_k^T \Delta \mathbf{w}_k|^2}{\sum_{i=1, i \neq k}^{K} p_i |\mathbf{h}_k^T \mathbf{w}_i|^2 + \sigma_n^2} \right)
\]

subject to

\[
\Delta_{n,n} \in \{0, 1\},
\]

\[
\sum_{n=1}^{N} \Delta_{n,n} = N_s,
\]

where \( \Delta \) is a binary-valued \( N \times N \) diagonal selection matrix. Providing the optimal solution requires an exhaustive search over \( \binom{N}{N_s} \) different combinations of antenna subsets, therefore it becomes prohibitive for large values of \( N \) and \( N_s \).

Accordingly, we propose in this paper two low-complexity, yet highly efficient, AS algorithms to solve the optimization problem in (10), and compare the performance of our proposed methods with the greedy selection algorithm proposed in [23, Algorithm 4]. It is noteworthy that the authors in [23] also considered a single-cell MU M-MIMO downlink system. However, their work focused on maximizing the sum rate through joint power allocation and AS as well as user scheduling, while in this work we focus only on the low-complexity design of AS schemes. It should be noted that throughout this work, we assume frequency-flat fading channels. Similar assumptions in M-MIMO systems with AS were made in [19], [26].
III. PROPOSED ANTENNA SELECTION ALGORITHMS

Although MF is one of the most attractive forms of linear precoding, it suffers from inter-user interference, which causes a performance degradation. Both of the proposed algorithms depend on the channel cross-correlation matrix $U = HH^H$, which can be expressed as

$$
U = \begin{pmatrix}
    h_1^T h_1 & h_1^T h_2 & \cdots & h_1^T h_K \\
    h_2^T h_1 & h_2^T h_2 & \cdots & h_2^T h_K \\
    \vdots & \vdots & \ddots & \vdots \\
    h_K^T h_1 & h_K^T h_2 & \cdots & h_K^T h_K \\
\end{pmatrix},
$$

where the elements on the diagonal of (11) are related to the desired signal gain for each user, while the elements in the upper and lower triangular parts are directly related to the inter-user interference, since the interference from user $j$ to user $i$ is $\frac{|h_i^* h_j|^2}{h_i^T h_i}$, for $i, j = 1,2,..,K$, and $i \neq j$. Each element in (11) is a summation of $N$ multiplications between two complex numbers, i.e.

$$
h_i^T h_j^* = h_i,1 h_j,1 + h_i,2 h_j,2 + \cdots + h_i,N h_j,N. \tag{12}
$$

Both of the proposed algorithms aim to maintain low implementation complexity by avoiding vector multiplications during the iterative selection process as explained in the following subsections.

A. Proposed User-Centric Antenna Selection (UCAS) Algorithm

In this method, the antennas at the BS are divided into $K$ groups based on their channel norms, where each group corresponds to one user. Moreover, each group should be allocated exactly $N/K$ antennas, therefore, any group that reaches its maximum limit will not be considered in the allocation process for the remaining set of available antennas. The $n$th antenna at the BS will be allocated to the $k$th group, denoted as $G_k$, only if it satisfies the following condition

$$
h_{k,n} \in G_k \iff |h_{k,n}| > |h_{i,n}|, \forall i \in \mathcal{K} \setminus k, \tag{13}
$$

where $\mathcal{K}$ is a set containing the indices of groups with less than $N/K$ antennas. For example, the 2nd antenna at the BS will be allocated to the 3rd group $G_3$, if and only if $|h_{3,2}| > |h_{i,2}|, \forall i = 1,..,K, i \neq 3$, and $G_3$ has less than $N/K$ antennas. The proposed grouping strategy is explained in detail in Algorithm 1.

At a given iteration, the UCAS algorithm aims to maximize the SINR for only the $k$th user, by selecting a single antenna from $G_k$. For example, let $H = [h_1, \ldots, h_K]^T \in \mathbb{C}^{K \times 1}$ be the channel matrix that corresponds to the $(t-1)$ selected antennas after as many iterations. Then, at the $t$th iteration, the antenna $\zeta$ will be selected as follows

$$
\zeta = \arg \max_{n \in G_k} \frac{h_k^T h_\zeta^* + h_{k,n} h_\zeta^*}{\sum_{i \neq k} |h_i^T h_\zeta^* + h_{i,n} h_\zeta^*|^2 + \sigma_n^2}. \tag{14}
$$

Algorithm 1 Proposed grouping strategy for the UCAS method

1. Input $K$, $N$, and $H$.
2. Initialize
3. $G_k = \emptyset, \forall k \in \{1, \ldots, K\}$
4. $t_k = 0, \forall k \in \{1, \ldots, K\}$ (Number of antennas in each group),
5. $\mathcal{K} = \{1, \ldots, K\}$ (Set of indices of groups with less than $N/K$ antennas),
6. for $n = 1 \to N$
7. $k^* = \arg \max_{k \in \mathcal{K}} |h_{n,k}|,$
8. $t_k^* = t_{k^*} + 1,$
9. $G_{k^*}(t_{k^*}) = n, (G_k(i) = (k)th\text{ element of } G_k),$
10. if $t_{k^*} = N/K$
11. $\mathcal{K} := \mathcal{K} \setminus k^*,$
12. end if
13. end for
14. Output $G_k, \forall k$

Moreover, we aim to achieve further complexity reduction by avoiding any vector multiplications in the numerator and denominator of (14). Accordingly, we store the values of the $N$ complex multiplications in (12), and for each element in (11), in a matrix $\Xi \in \mathbb{C}^{K \times N}$, which can be expressed as follows

$$
\Xi = \begin{pmatrix}
    \xi_{1,1} & \cdots & \cdots & \xi_{1,N} \\
    \vdots & \ddots & \vdots & \vdots \\
    \vdots & \vdots & \cdots & \vdots \\
    \xi_{K,1} & \cdots & \cdots & \xi_{K,N} \\
\end{pmatrix}, \tag{15}
$$

where $\xi_{i,1} = h_{1,i} h_{1,1}^* + h_{1,N} h_{1,N}^*, \xi_{i,N} = h_{i,1} h_{i,N}^*, \xi_{K,1} = h_{K,1} h_{K,K}^*$ and $\xi_{K,2,N} = h_{K,N} h_{K,N}^*$ (see Appendix A for details). Therefore, the SINR for the $k$th user depends mainly on rows $(K(k-1)+1, K(k-1) + 2, \ldots, K(k-1) + \sigma_n)$ in $\Xi$. Let $R_k$ denote the set of row indices in $\Xi$ that corresponds to user $k$, with $R_k$ being the $t$th element of $R_k$, for example, $R_1 = \{1, 2, \ldots, K\}$, while $R_{11} = 1$, and let $\omega = [\omega_1, \omega_2, \ldots, \omega_K]^T$ be a vector initialized with zeros, and used to update the values of the signal and interference terms after each antenna is selected, since selecting any antenna from any group will result in adding a column from $\Xi$ to $\omega$. To clarify this, assume at iteration $t$, the antenna $\zeta$ was selected, the vector $\omega$ then will be updated as follows

$$
\omega^{[t]} = \omega^{[t-1]} + \xi_{\zeta}, \tag{16}
$$

and the selection of $\zeta$ can now be expressed as follows

$$
\zeta = \arg \max_{n \in G_k} \frac{\left|\sum_{i=1}^{K} \omega_{R_k,i}^{[t-1]} + \xi_{R_k,n}^{[t]}\right|^2}{\sum_{i \neq k} \left|\sum_{i=1}^{K} \omega_{R_k,i}^{[t-1]} + \xi_{R_k,n}^{[t]}\right|^2 + \sigma_n^2}. \tag{17}
$$

Note that the selection process is now carried out without any vector multiplications. It should be noted that $G_k$ is updated after the selection, by removing the index that corresponds to the selected antenna before the next iteration starts. In the next iteration, the selection will be carried out in a similar manner to maximize the SINR for the next user, i.e. the $(k+1)$th
Algorithm 2 Proposed UCAS algorithm
1: Input $K$, $N$, $N_s$, $G_k(∀k)$, $R_k(∀k)$, and $H$
2: Initialize
3: $ω = 0_{K^2×1}$, $m = 0$, $Ξ = 0_{K^2×N}$
4: for $i = 1 → K$
5: for $j = 1 → K$
6: $m = m + 1$
7: for $n = 1 → N$
8: $ξ_{m,n} = h_{i,n}h_{j,n}^*$
9: end for
10: end for
11: end for
12: for $i = 1 → N_s/K$
13: for $k = 1 → K$
14: $ζ = \arg \max_n_{n_k}$\text{∈}$\mathcal{G}_k\{ \frac{∂ξ_{k,n} + ξ_{k,n}^*}{\sum_{i=1}^{K} ξ_{k,i} + ξ_{k,i}^* + σ_n^2} \}$
15: $ω := ω + ζ$
16: $\mathcal{G}_k := \mathcal{G}_k \setminus ζ$
17: end for
18: end for
19: Output $[H]_{ij} ∀ i ∈ S, ∀ k$

B. Proposed Semiblind Interference Rejection Antenna Selection (SIRAS) Algorithm

The SIRAS algorithm aims to minimize the interuser interference by minimizing the terms in the upper and lower triangle parts of (11) relying only on the signs of the interference terms, hence the term semiblind. Furthermore, since $h_i^T h_j = (h_i^T h_j)^*$, it is sufficient to minimize the elements in the upper triangular part only, and that will lead to exactly the same minimization for the terms in the lower triangular part, and vice versa. The algorithm starts by storing the $M × N$ complex multiplications in $Φ$, where $M = (K^2 - K)/2$ is the number of interference terms in the upper triangular part of (11), and it can be expressed as follows

$$Φ = \begin{pmatrix}
φ_{1,1} & \ldots & \ldots & φ_{1,N} \\
\vdots & \ddots & \ddots & \vdots \\
\vdots & \ddots & \ddots & \vdots \\
φ_{M,1} & \ldots & \ldots & φ_{M,N}
\end{pmatrix}. \quad (18)$$

Since the upper triangular part of (11) is considered in this work, the $N$ elements in the first row of (18) are the values of the $N$ complex number multiplications in the first interference term in (11), i.e. $φ_{1,n} = h_{1,n} h_{j,n}^*$, while $φ_{M,n} = h_{K-1,n} h_{K,n}$, for $n = 1, 2, \ldots, N$ (see Appendix A for details). The algorithm then selects its first antenna based on the maximum total channel norms

$$ζ^{[0]} = \arg \max_{n∈1,N} \|h_{n}^c\|, \quad (19)$$

where $ζ^{[0]}$ is the first selected antenna. After selecting the first antenna, a vector $ψ ∈ C^{M×1}$ will be initialized with the column in $Φ$ that corresponds to the selected antenna, i.e. $ψ^{[0]} = φ_{ζ^{[0]}}$. Note that the $M$ values in $ψ$ are complex numbers, and they are directly related to the interference between the users at any given iteration, hence minimizing these values will result in minimizing the total interference and therefore higher SINRs. Moreover, selecting any antenna will result in adding a column from $Φ$ to $ψ$. Assume that at iteration $t$, the antenna $ζ$ was selected, then $ψ$ will be updated as follows

$$ψ_{t} = ψ_{t-1} + φ_{ζ}^c, \quad (20)$$

therefore, the goal is to select the antenna that corresponds to the column in $Φ$ which will minimize the $M$ complex values in $ψ$. In other words, at iteration $t$, the algorithm aims to select the antenna $ζ$ from the set $S$, where the $n$th antenna (denoted as $λ_n$) belongs to $S$ if it satisfies the following condition

$$λ_n ∈ S ⇔ \left\{ \begin{array}{l}
\text{sign}(ℜ[φ_{m,n}]) \neq \text{sign}(ℜ[φ_{m,n}]) \\
\text{sign}(ℑ[φ_{m,n}]) \neq \text{sign}(ℑ[φ_{m,n}])
\end{array} \right\} \cap \left\{ \begin{array}{l}
\text{sign}(ℜ[φ_{m,n}]) \neq \text{sign}(ℜ[φ_{m,n}]) \\
\text{sign}(ℑ[φ_{m,n}]) \neq \text{sign}(ℑ[φ_{m,n}])
\end{array} \right\}, ∀m, n ∈ A, \quad (21)$$

where $A$ is a set containing the indices of available antennas at a given iteration. However, it is not guaranteed to find an antenna that satisfies the condition in (21), therefore, we relax this condition and select the antenna that has the highest number of opposite signs between $Φ$ and $ψ$. In the next iteration, the vector $ψ$ will be updated according to (20), and the same procedure will be repeated again until maximum
number of selected antennas is reached. The proposed method is described in detail in Algorithm 3.

IV. NUMERICAL RESULTS AND DISCUSSION

We apply Monte Carlo simulations to evaluate the spectral and energy efficiency performances of our proposed AS techniques for a wide range of Signal-to-Noise Ratios (SNRs). Furthermore, we compare our proposed algorithms with two low complexity AS methods. The first method is the Maximum SNR (MS) selection, where the $N_s$ antennas with the highest channel gains are activated to serve the users in the cell. While the second technique is the low complexity greedy method proposed in [23, Algorithm 4]. However, it is worth mentioning that the work in [23] considers joint power allocation and AS, while here we focus only on the AS part.

Before discussing the results, we need to introduce the parameters used in this work. First, we define the SNR per user as

$$\text{SNR} = \frac{\sigma_n^2}{\sigma_h^2} \frac{P_T}{K},$$

where the channel variance $\sigma_h^2$ can be found by using the path loss formula for a general urban channel model, which can be expressed as follows [27]

$$\text{PL (dB)} = 10 \log_{10} d^\alpha + \beta,$$

where $d$ is the distance between the users and the BS, $\alpha$ is the path loss component, and $\beta$ is the fixed-loss component. In our simulations, and unless stated otherwise, $d$ was set to 100 meters, $\alpha$ was assumed to be 2 as in [28], and $\beta = 10$ dB, therefore, $\sigma_h^2 = 10^{-5}$, while the noise variance $\sigma_n^2$ was assumed to have a value of $10^{-6}$. Our assumption for equal distances between users and the BS is based on the fact that different distances from the BS for different users can be tackled through power allocation techniques based on large scale fading [29]. In addition, the bandwidth $B$ was set to 20 MHz, the efficiency of the power amplifier $\eta$ was 0.35, while the receive and transmit circuit power consumption per RF chain $P_{rx}$ and $P_{tx}$ were set to 62.5 mW and 48.2 mW, respectively [25]. Finally, the simulations were averaged over $10^4$ different channel realizations for each SNR value.

A. Impact of AS on spectral and energy efficiencies

The advantage of designing efficient AS schemes when MF precoding is applied, in terms of spectral and energy efficiencies, can be seen in Figs. 2 and 3, respectively. Both of the proposed algorithms achieve higher rates with better EE than the full system case, where all the antennas are employed, i.e. when $N_s = N$. Furthermore, in terms of sum rate performance, increasing the SNR results in decreasing the number of antennas required to match the achievable rate where all antennas are activated. The reason behind this is that as the SNR increases, the effect of the noise becomes negligible, and the main factor that degrades the performance is the inter-user interference, and our proposed algorithms select the antennas that highly minimize the inter-user interference. In other words, although selecting a subset of the available antennas will reduce the signal gain, the resultant SNR will be higher, hence higher rates are achieved. Moreover, in terms of EE, applying AS techniques can dramatically improve the system performance as can be seen in Fig. 3. For example, when the number of selected antennas is 64, the UCAS and SIRAS algorithms outperform the full system case by 190.8 and 184 Mbits/Joule, respectively.

B. Achievable rates with fixed number of selected antennas

In this subsection, we show the performance of the proposed algorithms in terms of the achievable rates, and compare our results with MS and greedy selection techniques for a wide range of SNR values.

Fig. 4 shows the total sum rate when the BS is equipped with 128 antennas, out of which 32 are selected, for different AS schemes. The proposed algorithms outperform significantly both greedy and MS methods. Moreover, UCAS slightly outperforms the SIRAS algorithm when the SNR is less than 10 dB, while they both show the same performance for higher SNR values. For this scenario, employing all the available antennas at the BS achieves higher rates than our proposed algorithms when the number of selected antennas is 32.

Fig. 5 shows the total sum rate when 256 antennas are placed at the BS, and 64 antennas are selected by the proposed algorithms. The proposed methods not only outperform the greedy and MS techniques, but also the case where all the antennas are activated for SNR values higher than 0.5 dB and 3.5 dB for the UCAS and SIRAS algorithms, respectively.

From Fig. 5, the UCAS algorithm outperforms the full system, greedy, and MS selection methods at 10 dB SNR by 5.44, 10.12, and 19.64 bps/Hz, respectively, which correspond to a significant 108.8, 202.4, and 392.8 Mbps for a system with 20 MHz of bandwidth, while the SIRAS method shows an improvement of 2.46, 7.14, and 16.66 bps/Hz compared to the full system, greedy, and MS selection methods, respectively.
C. Energy Efficiency Performance

In this subsection, we show the EE performance for the different AS methods as well as the full system case. We assume 128 antennas are placed at the BS, out of which 32 are selected to serve 8 users in the cell.

As Fig. 6 shows, at moderate SNR values, the proposed methods outperform significantly all other schemes for both scenarios. Furthermore, activating all the antennas results in extremely poor EE performance, which validates the importance of AS for energy efficient systems. For example, the proposed methods outperform the greedy algorithm, MS selection, and the full system by 59 Mbits/Joule, 92 Mbits/Joule, and 194 Mbits/Joule, respectively, at SNR of 10 dB.

D. Achievable rates with imperfect channel state information

In practical scenarios, the complex channel matrix $H$ is estimated at the BS through pilot signals sent by the users, and channel estimation errors arise in any practical system. Different techniques have been designed to obtain the CSI in both conventional and M-MIMO systems [30]–[32]. In general, increasing the number of pilot signals leads to a better
estimation accuracy, but that comes at the price of a reduced data transmission interval. However, for M-MIMO systems, it was shown that a minimum of $K$ pilot signals per coherence interval are enough when the system is operating in a TDD mode [33]. In this subsection, we evaluate the performance of the proposed algorithms under imperfect CSI. The estimated channel matrix $\hat{\mathbf{H}}$ can be given as [34], [35]

$$\hat{\mathbf{H}} = \mathbf{H} + \epsilon \mathbf{E}, \quad (24)$$

where $\epsilon \mathbf{E}$ represents the channel estimation error term, and is uncorrelated with $\mathbf{H}$. The entries of $\mathbf{E}$ are independent and identically distributed random variables with zero mean and variance of $\sigma^2_h$. Furthermore, $\epsilon$ controls the estimation accuracy, and $\epsilon = 0$ means the BS has a perfect CSI.

Fig. 7 demonstrates the achievable rates of the proposed methods for different levels of channel estimation error. The achievable rates are effected by the channel estimation accuracy for both methods. In addition, at higher SNR, the proposed algorithms become more sensitive under imperfect CSI. The reason behind this is that as the SNR increases, the noise effect becomes negligible, and the presence of CSI errors leads to the selection of sub-optimal antennas. For example, when the SNR $= -5 \text{ dB}$, the UCAS and SIRAS algorithms suffer a 3.57 and 3.23 bps/Hz rate loss when $\epsilon = 0.4$ compared to the perfect CSI case, respectively; while the achievable rates for the same algorithms degrade by 7.92 and 7.11 bps/Hz, respectively, when $\epsilon = 0.4$, at SNR of 10 dB compared to the perfect CSI case.

Fig. 8 demonstrates the performance of different AS schemes for a wide range of SNR values when $\epsilon = 0.2$. The UCAS and SIRAS algorithms outperform the full system case at 2 dB and 7 dB, respectively, compared to 0 dB and 4 dB for perfect CSI. Furthermore, at SNR $= 10 \text{ dB}$, the UCAS algorithm outperforms the greedy method and MS selection by 8.34 bps/Hz and 17.34 bps/Hz, respectively; while the SIRAS method achieves 5.91 bps/Hz and 14.91 bps/Hz higher rates than the greedy method and MS selection, respectively, at the same SNR value.

E. Performance of the proposed methods with multiple random initializations

The performance of both of the proposed methods can be further enhanced by adopting multiple random initialization-\hspace{1pt}s of users and antennas, for UCAS and SIRAS schemes, respectively. More specifically, for the UCAS method, after obtaining $\Xi$ and grouping the antennas into $K$ groups, the users were ordered from 1 to $K$, where for each of the $N_s/K$ rounds of selection, one antenna was selected to maximize the SINR for the 1st user, followed by selecting one antenna to maximize the SINR for the 2nd user, and so forth until the $K$th user. However, adopting multiple random initialization orders of the $K$ users can further improve the performance by obtaining multiple sets of selected antennas, and then choose the set that achieves the highest SINR. Similarly, and for the SIRAS scheme, multiple random antennas can be chosen as the first selected antenna, and then perform the selection for the following $N_s - 1$ iterations to obtain different sets of selected antennas, and finally select the set that achieves the maximum total SINR.

As demonstrated in Fig. 9, both methods achieve a considerable gain when multiple random initial ordering of users and antennas are considered. Specifically, at a SNR of 10 dB, the UCAS scheme shows an improvement of 3.73 bps/Hz when 10 random initial orders ($L = 10$) of users are considered compared to the classical single ordering case ($L = 1$) shown in Algorithm 2. Moreover, the SIRAS scheme achieves a 4.03 bps/Hz gain when considering 10 random antennas as the first selected antenna ($L = 10$), compared to the case where the first selected antenna is only the one with the highest channel gain ($L = 1$). However, although these methods can provide
a considerable performance gain, they come at the price of increased complexity for their implementations as will be discussed in Section V.

**F. Performance of UCAS algorithm under different grouping strategies**

Throughout this work, grouping the antennas into $K$ groups for the UCAS scheme was performed according to our proposed method in Algorithm 1. However, there are different methods to perform this part of the algorithm, and in this section, we will investigate the performance of two additional grouping strategies.

The first method is the random grouping strategy, where each group will randomly be assigned $N/K$ antennas. A more efficient way of grouping is assigning the antenna that has the maximum channel gain with only the first user as the first assigned antenna for the first group. Followed by assigning the antenna that has the maximum channel gain with only the second user to the second group, and so forth until the $K$th group. Moreover, the same procedure will be repeated until each group has $N/K$ antennas assigned to it. We call the latter scheme as the user-based grouping.

As demonstrated in Fig. 10, the random grouping method shows an extremely poor performance compared to the other two grouping schemes. In contrast, the user-based grouping method shows a marginal sum rate improvement compared to the original grouping strategy given in Algorithm 1. However, the user-based grouping method requires higher complexity as will be demonstrated in Section V. Therefore, the original grouping method demonstrates a positive performance-complexity trade-off compared to the other two schemes.

**G. Fairness measurement of the proposed AS schemes**

To have more insight into the performance of the proposed AS schemes, we measure the equality among different users by applying Jain’s equality formula, which can be defined as follows [36]

\[ \text{Jain’s fairness index} = \frac{\left( \sum_{k=1}^{K} R_k \right)^2}{K \sum_{k=1}^{K} R_k^2}. \]  

As shown in Fig. 11, the proposed methods achieve near optimal fairness among the users, as the fairness index exceeds 97% for a wide range of SNR values. Moreover, the UCAS method, and regardless of the grouping strategy, can deliver higher fairness to users than the SIRAS scheme for the entire range of SNR values considered in this work.
TABLE II
NUMBER OF ADDITIONS AND MULTIPLICATIONS FOR DIFFERENT AS SCHEMES

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Operator</th>
<th>Additions</th>
<th>Multiplications</th>
</tr>
</thead>
<tbody>
<tr>
<td>MS</td>
<td>$|h_{m}^{\nu}|$, $\forall \nu$</td>
<td>$N(2K - 1)$</td>
<td>$N(2K + 1)$</td>
</tr>
<tr>
<td></td>
<td>Sorting</td>
<td>$N \log N$</td>
<td>$-$</td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>$N(2K + \log N - 1)$</td>
<td>$N(2K + 1)$</td>
</tr>
<tr>
<td>Greedy</td>
<td>$|h_{m}^{\nu}|$, $\forall \nu$</td>
<td>$N(2K - 1)$</td>
<td>$N(2K + 1)$</td>
</tr>
<tr>
<td></td>
<td>$w_{k}, \forall k$</td>
<td>$\sum_{l=1}^{N} K(2l - 1)$</td>
<td>$\sum_{l=1}^{N} K(4l + 1)$</td>
</tr>
<tr>
<td></td>
<td>$\gamma_{k}, \forall k$</td>
<td>$K \sum_{l=1}^{N} (4Kl - 1)$</td>
<td>$K \sum_{l=1}^{N} (4Kl + 5K + 1)$</td>
</tr>
<tr>
<td></td>
<td>$R$</td>
<td>$(2K - 1)N_{s}$</td>
<td>$-$</td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>$N(2K + \log N - 1) - N_{s} + \sum_{l=1}^{N} l(4K^{2} + 2K)$</td>
<td>$N(2K + 1) + N_{s}(5K^{2} + 2K) + \sum_{l=1}^{N} l(4K^{2} + 4K)$</td>
</tr>
<tr>
<td>UCAS</td>
<td>$\Xi$</td>
<td>$2K^{2}N$</td>
<td>$4K^{2}N$</td>
</tr>
<tr>
<td></td>
<td>$G_{k}, \forall k$</td>
<td>$2KN$</td>
<td>$3KN$</td>
</tr>
<tr>
<td></td>
<td>$\zeta_{k}, \forall {1, \ldots, N_{s}}$</td>
<td>$\sum_{l=1}^{N_{s}}/K(N - l)4K^{2}$</td>
<td>$\sum_{l=1}^{N_{s}}/K(N - l)(4K^{2} + K)$</td>
</tr>
<tr>
<td></td>
<td>$\omega_{k}, \forall {1, \ldots, N_{s}}$</td>
<td>$2K^{2}N_{s}$</td>
<td>$-$</td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>$K(2KN + 2N + 2KN_{s}) + \sum_{l=1}^{N_{s}}/K(N - l)4K^{2}$</td>
<td>$N(4K^{2} + 3K) + \sum_{l=1}^{N_{s}}/K(N - l)(4K^{2} + K)$</td>
</tr>
<tr>
<td>SIRAS</td>
<td>$\Phi$</td>
<td>$N(2K^{2} - K)$</td>
<td>$2N(2K^{2} - K)$</td>
</tr>
<tr>
<td></td>
<td>$|h_{m}^{\nu}|$, $\forall \nu$</td>
<td>$N(2K - 1)$</td>
<td>$N(2K + 1)$</td>
</tr>
<tr>
<td></td>
<td>$\zeta_{k}$</td>
<td>$N$</td>
<td>$-$</td>
</tr>
<tr>
<td></td>
<td>$\lambda$</td>
<td>$\sum_{l=1}^{N_{s}-1} 2(N - l)(K^{2} - K)$</td>
<td>$-$</td>
</tr>
<tr>
<td></td>
<td>$\zeta_{1, \ldots, N_{s}-1}$</td>
<td>$\sum_{l=1}^{N_{s}-1} (N - l)$</td>
<td>$-$</td>
</tr>
<tr>
<td></td>
<td>$\psi_{1, \ldots, N_{s}-1}$</td>
<td>$(N_{s} - 1)(K^{2} - K)$</td>
<td>$-$</td>
</tr>
<tr>
<td>Total</td>
<td></td>
<td>$N(K^{2} + K) + (N_{s} - 1)(K^{2} - K) + \sum_{l=1}^{N_{s}-1} (N - l)(2K^{2} - 2K + 1)$</td>
<td>$N(2K^{2} + 1)$</td>
</tr>
</tbody>
</table>

V. COMPLEXITY ANALYSIS

In TDD M-MIMO systems with AS, the BS is required for each coherence interval, to perform channel estimation and AS to enable efficient uplink and downlink data transmission. Therefore, the complexity analysis presented in this section is of great importance to evaluate the efficiency of the proposed AS schemes in practical scenarios, since higher complexity requires longer processing time, thus reducing the data transmission interval, and vice versa. It is worth highlighting that state-of-the-art Digital Signal Processors (DSPs), such as that found on the TCI6638K2K DSP board designed by Texas instruments for the next generation of wireless networks, can support up to 153.6 GigaFLOPs per second [37].

We follow the complexity analysis in [38], where the addition between two real numbers is equivalent to 1 FLOP, while their multiplication is equivalent to 4 additions. Moreover, we assume that comparing the signs of two real numbers is equivalent to 1 addition, while finding the square root of a real number and division between two real numbers are each equivalent to one real multiplication.

A. Complexity of MS scheme

The MS method has the lowest complexity, however that comes at the price of performance degradation. Finding the channel norms for the $N$ antennas requires $N(10K + 3)$ FLOPs, while sorting the antennas requires $N \log_{10} N$ FLOPs, therefore, the total complexity is

$$C_{MS} = N(10K + \log_{10} N + 3).$$

B. Complexity of greedy AS algorithm [23, Algorithm 4]

The greedy algorithm starts by finding the channel norms for the $N$ antennas and sorting them, which requires $N(10K + \log_{10} N + 3)$ FLOPs. The algorithm will then start its iterations, where at each iteration the sum rate in (7) will be evaluated. Finding the MF precoding weights for all the users requires $K \sum_{l=1}^{N_{s}} (18l + 3)$ FLOPs, while evaluating the SINRs for all users requires $\sum_{l=1}^{N_{s}} (20K^{2}l + 20K^{2} + 3K)$ FLOPs. Finally, calculating the sum rate after obtaining the SINR values takes $(2K - 1)N_{s}$ FLOPs, therefore, the total complexity of the greedy algorithm can be given as

$$C_{greedy} = N(10K + \log_{10} N + 3) + N_{s}(20K^{2} + 8K - 1) + \sum_{l=1}^{N_{s}} l(20K^{2} + 18K).$$

However, here we only consider an ideal case for the greedy algorithm by assuming that each selected antenna will boost
the total sum rate, therefore, the actual complexity of the greedy algorithm is higher than that shown in this section, otherwise it would have shown the exact same sum rate performance as the MS scheme.

C. Complexity of UCAS algorithm with a single initial order of users (L = 1)

The UCAS algorithm starts by evaluating the matrix $\Xi$ that has $K^2 \times N$ complex multiplications, which results in $18K^2N$ FLOPs. Then, the absolute channel values between the $N$ antennas and $K$ users are evaluated before being allocated to $\psi_k, \forall k$ according to Algorithm 1, which results in $14KN$ FLOPs. The $N_s/K$ iterations will then start to select $K$ antennas at each iteration, which results in $\sum_{s=1}^{N_s/K}(20K + 4)(N - K)$ FLOPs. Finally, the vector $\omega$ is updated after selecting each antenna, which requires $2K^2N_s$ FLOPs. Therefore, the total number of FLOPs required by the UCAS algorithm can be given as

$$C_{\text{UCAS}} = KN(18K + 14) + 2K^2N_s$$

$$+ \sum_{l=1}^{N_s/K} (20K + 4)(N - Kl).$$  \hspace{1cm} (28)

D. Complexity of SIRAS algorithm with a single initial order of antennas (L = 1)

The SIRAS algorithm requires $9N(K^2 - K)$ FLOPs to obtain $\Phi$, after that, $N(10K + 4)$ FLOPs are required to select the first antenna. The $N_s - 1$ iterations will then start, and at each iteration the vector $\lambda$ is used to store the number of opposite signs, which takes $\sum_{l=1}^{N_s-1} 2(N - l)(K^2 - K)$ FLOPs. Finding the $N_s - 1$ antennas from $\lambda$ requires $\sum_{l=1}^{N_s-1} (N - l)$ FLOPs and finally, updating the vector $\psi$ results in $(N_s - 1)(K^2 - K)$ FLOPs. Therefore, the total complexity for the SIRAS algorithm can be given as

$$C_{\text{SIRAS}} = N(9K^2 + K + 4) + (K^2 - K)(N_s - 1)$$

$$+ \sum_{l=1}^{N_s-1} (N - l)(2K^2 - 2K + 1).$$ \hspace{1cm} (29)

The numbers of additions and multiplications for the different AS schemes are shown for each operator in Table II\(^1\). It should be noted that in the aforementioned table, the number of additions also includes the comparison between the values or signs of two real numbers, while the number of multiplications includes the division as well as finding the square root of a real number.

Figs. 12 and 13 show the complexity for different AS schemes for different numbers of antennas $N$ at the BS and different numbers of users $K$, respectively, with $N_s = N/4$. In both results, the MS algorithm has the lowest complexity, followed by the UCAS algorithm, then the SIRAS, while the greedy algorithm requires the highest complexity among the adopted schemes. As mentioned before, the actual complexity for the greedy algorithm is higher than that shown in this section, and here we only consider an ideal case.

To gain more insight on the efficiency of our proposed schemes in practical scenarios, assume that we have a BS equipped with 256 antennas, out of which 64 are selected to serve 8 single-antenna users. According to the analysis presented in this section, both of the proposed methods require less than $2 \times 10^6$ FLOPs, which corresponds to 13 $\mu$s of processing time using the aforementioned DSP board. Note that this is an impressive outcome for the proposed methods given that in the current Long-Term Evolution (LTE) standard, transmitting one symbol of data requires 71.4 $\mu$s [11].

\(^1\)The number of additions and multiplications of UCAS and SIRAS methods shown in Table II are for the case where $L = 1$. 
E. Complexity analysis of UCAS algorithm with multiple random ordering of users \((L > 1)\)

The complexity required for obtaining \(L\) different candidate sets of selected antennas for the UCAS method, with \(L > 1\), can be given as follows

\[
C_{\text{UCAS}}^{[L]} = L C_{\text{UCAS}}^{[1]} - (L - 1)(18K^2 N + 14KN) + C_{\text{add}}^{[L]},
\]

where \(C_{\text{UCAS}}^{[1]}\) is the complexity of UCAS method for a single initialization given in (28), and \(C_{\text{add}}^{[L]}\) is the additional complexity required to evaluate the SINRs for the \(L\) candidates of different antenna subsets and selecting the one that has the maximum SINR as the final solution. For any number of random initializations \(L > 1\), \(C_{\text{add}}^{[L]}\) can be given as follows

\[
C_{\text{add}}^{[L]} = L(20K^2N_s + 20K^2 + 18KN_s + 7K).
\]

Finally, the second term in the right hand side of (30) is subtracted from the total complexity due to the fact that the matrix \(\Xi\) has to be evaluated only once, as well as grouping the antennas into the \(K\) groups according to Algorithm 1.

F. Complexity of SIRAS method with multiple random initializations of antennas \((L > 1)\)

Similar to the UCAS method, the complexity of the SIRAS algorithm when \(L > 1\) can be given as follows

\[
C_{\text{SIRAS}}^{[L]} = L C_{\text{SIRAS}}^{[1]} - (L - 1)(9NK^2 + NK + 4N) + C_{\text{add}}^{[L]},
\]

where \(C_{\text{SIRAS}}^{[1]}\) is the complexity of the SIRAS method for a single initialization given in (29), \(C_{\text{add}}^{[L]}\) is the additional complexity to compute the SINRs of the \(L\) different antenna subsets given in (31), while the second term of the right hand side of (32) refers to the fact that the matrix \(\Phi\), sorting the antennas, and selecting the first antenna as the one with the maximum channel gain needs to be carried out only once.

As shown in Fig. 14, both methods suffer from increased complexity as the number of random initial orders \((L)\) increases, with the SIRAS method requires higher complexity compared to the UCAS scheme.

G. Complexity of user-based grouping scheme

The complexity of the user-based grouping strategy requires \(\sum_{l=1}^{N} 14(N - l + 1)\) FLOPs to perform the grouping process, compared to \(14KN\) FLOPs for the original approach given in Algorithm 1. As demonstrated in Fig. 15, the user-based grouping scheme requires much higher complexity than the original proposed grouping method, especially as the number of antennas at the BS increases.

VI. CONCLUSIONS

In this paper, two efficient low complexity AS algorithms were proposed for downlink MU M-MIMO systems with MF precoding. Dramatic complexity reduction was achieved by storing the multiplications between each two entries of the channel matrix to avoid any vector multiplications during the iterative selection process. The first algorithm divides the available antennas into \(K\) groups depending on their channel norms, where each group corresponds to one user only. This resulted in a search space reduction by a factor of \(K\). The second algorithm was designed to reject the interuser interference relying only on the signs of the interference terms. The performance of the proposed methods under perfect and imperfect CSI was evaluated in terms of spectral and energy efficiencies. The proposed methods demonstrated great performance-complexity trade-off compared to other AS
methods as well as to the case where all the antennas were employed.

**APPENDIX A**

The entries of matrices $\Xi$ and $\Phi$, for a system with 4 single-antenna users and 8 antennas at the BS, i.e. $K = 4$ and $N = 8$, can be defined as shown in (33) and (34), respectively.

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**REFERENCES**


\[
\phi = \begin{bmatrix}
\xi_{1,1} & \xi_{1,2} & \xi_{1,3} & \xi_{1,4} & \xi_{1,5} & \xi_{1,6} & \xi_{1,7} & \xi_{1,8} \\
\xi_{2,1} & \xi_{2,2} & \xi_{2,3} & \xi_{2,4} & \xi_{2,5} & \xi_{2,6} & \xi_{2,7} & \xi_{2,8} \\
\xi_{3,1} & \xi_{3,2} & \xi_{3,3} & \xi_{3,4} & \xi_{3,5} & \xi_{3,6} & \xi_{3,7} & \xi_{3,8} \\
\xi_{4,1} & \xi_{4,2} & \xi_{4,3} & \xi_{4,4} & \xi_{4,5} & \xi_{4,6} & \xi_{4,7} & \xi_{4,8} \\
\xi_{5,1} & \xi_{5,2} & \xi_{5,3} & \xi_{5,4} & \xi_{5,5} & \xi_{5,6} & \xi_{5,7} & \xi_{5,8} \\
\xi_{6,1} & \xi_{6,2} & \xi_{6,3} & \xi_{6,4} & \xi_{6,5} & \xi_{6,6} & \xi_{6,7} & \xi_{6,8} \\
\xi_{7,1} & \xi_{7,2} & \xi_{7,3} & \xi_{7,4} & \xi_{7,5} & \xi_{7,6} & \xi_{7,7} & \xi_{7,8} \\
\xi_{8,1} & \xi_{8,2} & \xi_{8,3} & \xi_{8,4} & \xi_{8,5} & \xi_{8,6} & \xi_{8,7} & \xi_{8,8} \\
\xi_{9,1} & \xi_{9,2} & \xi_{9,3} & \xi_{9,4} & \xi_{9,5} & \xi_{9,6} & \xi_{9,7} & \xi_{9,8} \\
\xi_{10,1} & \xi_{10,2} & \xi_{10,3} & \xi_{10,4} & \xi_{10,5} & \xi_{10,6} & \xi_{10,7} & \xi_{10,8} \\
\xi_{11,1} & \xi_{11,2} & \xi_{11,3} & \xi_{11,4} & \xi_{11,5} & \xi_{11,6} & \xi_{11,7} & \xi_{11,8} \\
\xi_{12,1} & \xi_{12,2} & \xi_{12,3} & \xi_{12,4} & \xi_{12,5} & \xi_{12,6} & \xi_{12,7} & \xi_{12,8} \\
\xi_{13,1} & \xi_{13,2} & \xi_{13,3} & \xi_{13,4} & \xi_{13,5} & \xi_{13,6} & \xi_{13,7} & \xi_{13,8} \\
\xi_{14,1} & \xi_{14,2} & \xi_{14,3} & \xi_{14,4} & \xi_{14,5} & \xi_{14,6} & \xi_{14,7} & \xi_{14,8} \\
\xi_{15,1} & \xi_{15,2} & \xi_{15,3} & \xi_{15,4} & \xi_{15,5} & \xi_{15,6} & \xi_{15,7} & \xi_{15,8} \\
\xi_{16,1} & \xi_{16,2} & \xi_{16,3} & \xi_{16,4} & \xi_{16,5} & \xi_{16,6} & \xi_{16,7} & \xi_{16,8}
\end{bmatrix}
\]

\[
\Phi = \begin{bmatrix}
\phi_{1,1} & \phi_{1,2} & \phi_{1,3} & \phi_{1,4} & \phi_{1,5} & \phi_{1,6} & \phi_{1,7} & \phi_{1,8} \\
\phi_{2,1} & \phi_{2,2} & \phi_{2,3} & \phi_{2,4} & \phi_{2,5} & \phi_{2,6} & \phi_{2,7} & \phi_{2,8} \\
\phi_{3,1} & \phi_{3,2} & \phi_{3,3} & \phi_{3,4} & \phi_{3,5} & \phi_{3,6} & \phi_{3,7} & \phi_{3,8} \\
\phi_{4,1} & \phi_{4,2} & \phi_{4,3} & \phi_{4,4} & \phi_{4,5} & \phi_{4,6} & \phi_{4,7} & \phi_{4,8} \\
\phi_{5,1} & \phi_{5,2} & \phi_{5,3} & \phi_{5,4} & \phi_{5,5} & \phi_{5,6} & \phi_{5,7} & \phi_{5,8} \\
\phi_{6,1} & \phi_{6,2} & \phi_{6,3} & \phi_{6,4} & \phi_{6,5} & \phi_{6,6} & \phi_{6,7} & \phi_{6,8}
\end{bmatrix}
\]

\[
\Omega = \begin{bmatrix}
\h_1 h_1^* & \h_2 h_1^* & \h_3 h_1^* & \h_4 h_1^* & \h_5 h_1^* & \h_6 h_1^* & \h_7 h_1^* & \h_8 h_1^* \\
\h_1 h_2^* & \h_2 h_2^* & \h_3 h_2^* & \h_4 h_2^* & \h_5 h_2^* & \h_6 h_2^* & \h_7 h_2^* & \h_8 h_2^* \\
\h_1 h_3^* & \h_2 h_3^* & \h_3 h_3^* & \h_4 h_3^* & \h_5 h_3^* & \h_6 h_3^* & \h_7 h_3^* & \h_8 h_3^* \\
\h_1 h_4^* & \h_2 h_4^* & \h_3 h_4^* & \h_4 h_4^* & \h_5 h_4^* & \h_6 h_4^* & \h_7 h_4^* & \h_8 h_4^* \\
\h_1 h_5^* & \h_2 h_5^* & \h_3 h_5^* & \h_4 h_5^* & \h_5 h_5^* & \h_6 h_5^* & \h_7 h_5^* & \h_8 h_5^* \\
\h_1 h_6^* & \h_2 h_6^* & \h_3 h_6^* & \h_4 h_6^* & \h_5 h_6^* & \h_6 h_6^* & \h_7 h_6^* & \h_8 h_6^* \\
\h_1 h_7^* & \h_2 h_7^* & \h_3 h_7^* & \h_4 h_7^* & \h_5 h_7^* & \h_6 h_7^* & \h_7 h_7^* & \h_8 h_7^* \\
\h_1 h_8^* & \h_2 h_8^* & \h_3 h_8^* & \h_4 h_8^* & \h_5 h_8^* & \h_6 h_8^* & \h_7 h_8^* & \h_8 h_8^*
\end{bmatrix}
\]

\[
\Phi = \begin{bmatrix}
\h_1 h_1^* & \h_2 h_1^* & \h_3 h_1^* & \h_4 h_1^* & \h_5 h_1^* & \h_6 h_1^* & \h_7 h_1^* & \h_8 h_1^* \\
\h_1 h_2^* & \h_2 h_2^* & \h_3 h_2^* & \h_4 h_2^* & \h_5 h_2^* & \h_6 h_2^* & \h_7 h_2^* & \h_8 h_2^* \\
\h_1 h_3^* & \h_2 h_3^* & \h_3 h_3^* & \h_4 h_3^* & \h_5 h_3^* & \h_6 h_3^* & \h_7 h_3^* & \h_8 h_3^* \\
\h_1 h_4^* & \h_2 h_4^* & \h_3 h_4^* & \h_4 h_4^* & \h_5 h_4^* & \h_6 h_4^* & \h_7 h_4^* & \h_8 h_4^* \\
\h_1 h_5^* & \h_2 h_5^* & \h_3 h_5^* & \h_4 h_5^* & \h_5 h_5^* & \h_6 h_5^* & \h_7 h_5^* & \h_8 h_5^* \\
\h_1 h_6^* & \h_2 h_6^* & \h_3 h_6^* & \h_4 h_6^* & \h_5 h_6^* & \h_6 h_6^* & \h_7 h_6^* & \h_8 h_6^* \\
\h_1 h_7^* & \h_2 h_7^* & \h_3 h_7^* & \h_4 h_7^* & \h_5 h_7^* & \h_6 h_7^* & \h_7 h_7^* & \h_8 h_7^* \\
\h_1 h_8^* & \h_2 h_8^* & \h_3 h_8^* & \h_4 h_8^* & \h_5 h_8^* & \h_6 h_8^* & \h_7 h_8^* & \h_8 h_8^*
\end{bmatrix}
\]
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